

ESDP Relaxation of Sensor Network Localization Analysis, Extensions and Algorithm

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(Joint work with Paul Tseng)

Talk Outline

- Sensor network localization.
- ESDP relaxation: properties and accuracy certificate.
- A robust version of ESDP for the noisy case.
- LPCGD algorithm and numerical simulations.
- Conclusion and extension.

Sensor Network Localization

Basic Problem:

- n pts $\underbrace{x_1, \dots, x_m}_{\text{sensors}}, \underbrace{x_{m+1}, \dots, x_n}_{\text{anchors}}$ in \mathbb{R}^2 .
- Know last $n - m$ pts ('anchors') x_{m+1}, \dots, x_n and Eucl. dist. estimate for some pairs of 'neighboring' pts (i.e. within 'radio range')

$$d_{ij} \geq 0 \quad \forall (i, j) \in \mathcal{A},$$

with $\mathcal{A} \subseteq \{(i, j) : 1 \leq i < j \leq n\}$.

- Estimate the first m pts ('sensors') x_1, \dots, x_m .

Optimization Problem Formulation

$$v_p := \min_{x_1, \dots, x_m} \sum_{(i,j) \in \mathcal{A}} \left| \|x_i - x_j\|^2 - d_{ij}^2 \right|.$$

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- Objective function is nonconvex. m can be large ($m > 1000$).
- Problem is NP-hard (reduction from PARTITION).
- Use a convex relaxation.
- Fast, distributed algorithm.

SDP Relaxation

Let $X = [x_1 \cdots x_m]$. $Y = X^T X \iff Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \succeq 0, \text{rank } Z = 2.$

SDP Relaxation

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SDP relaxation (Biswas, Ye '03):

$$v_{\text{sdp}} := \min_Z \sum_{(i,j) \in \mathcal{A}^a} |y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2|$$

$$+ \sum_{(i,j) \in \mathcal{A}^s} |y_{ii} - 2y_{ij} + y_{jj} - d_{ij}^2|$$

$$\text{s.t. } Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \succeq 0.$$

Adding nonconvex constraint $\text{rank } Z = 2$ yields the original problem.

ESDP Relaxation

ESDP relaxation (Wang, Zheng, Boyd, Ye '07):

$$\begin{aligned}
 v_{\text{esdp}} := & \min_Z \sum_{(i,j) \in \mathcal{A}^a} |y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2| \\
 & + \sum_{(i,j) \in \mathcal{A}^s} |y_{ii} - 2y_{ij} + y_{jj} - d_{ij}^2| \\
 \text{s.t. } & Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \\
 & \begin{bmatrix} y_{ii} & y_{ij} & x_i^T \\ y_{ij} & y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \succeq 0 \quad \forall (i,j) \in \mathcal{A}^s \\
 & \begin{bmatrix} y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \succeq 0 \quad \forall i = 1, \dots, m.
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 \end{aligned}$$

In simulations, ESDP is nearly as strong as SDP relaxation, and solvable much faster by IP method.

Properties of ESDP

Assume

$$d_{ij} = \|x_i^{\text{true}} - x_j^{\text{true}}\| \quad \forall (i, j) \in \mathcal{A}.$$

“noiseless case”

$$(x_j^{\text{true}} = x_j \quad \forall j > m)$$

Fact 0:

$$Z^{\text{true}} := [X^{\text{true}} \quad I]^T [X^{\text{true}} \quad I] = \begin{bmatrix} (X^{\text{true}})^T X^{\text{true}} & (X^{\text{true}})^T \\ X^{\text{true}} & I \end{bmatrix}$$

is a soln of ESDP (i.e., $Z^{\text{true}} \in \text{Sol}(\text{ESDP})$).

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$$x_i \text{ is invariant over } \text{Sol}(\text{ESDP}) \Rightarrow x_i = x_i^{\text{true}}.$$

Let $\text{tr}_i[Z] := y_{ii} - \|x_i\|^2, \quad i = 1, \dots, m.$ “*i*th individual trace”

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Fact 2 (P, Tseng '09): For each i ,

$$x_i \text{ is invariant over } \text{Sol}(\text{ESDP}) \implies \text{tr}_i[Z] = 0 \forall Z \in \text{Sol}(\text{ESDP}).$$

In practice, there are measurement noises:

$$d_{ij}^2 = \|x_i^{\text{true}} - x_j^{\text{true}}\|^2 + \delta_{ij} \quad \forall (i, j) \in \mathcal{A}.$$

When $\delta := (\delta_{ij})_{(i,j) \in \mathcal{A}} \approx 0$, does $\text{tr}_i[Z] = 0$ (with $Z \in \text{ri}(\text{Sol}(\text{ESDP}))$) imply x_i is near the true position of sensor i ?

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No.

Fact 3 (P, Tseng '09): For $|\delta_{ij}| \approx 0$,

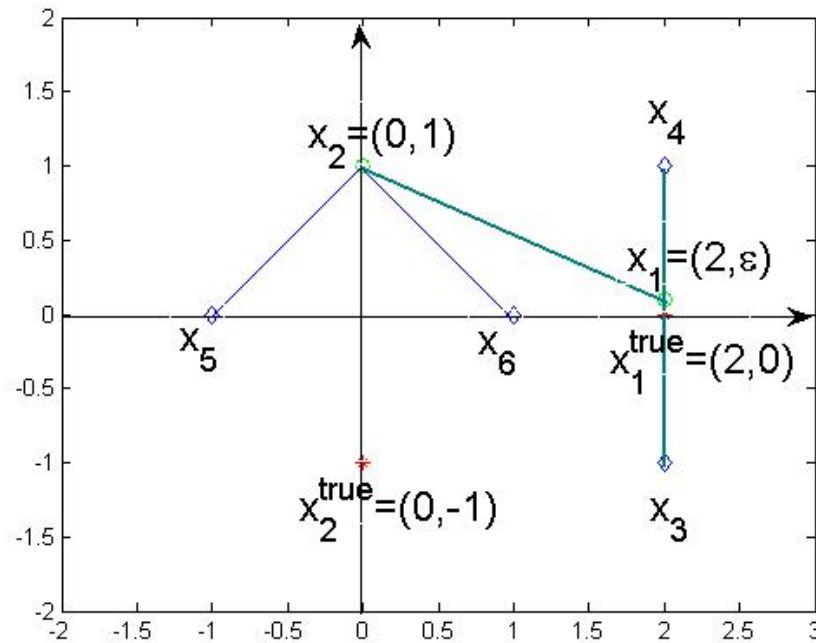
$$\text{tr}_i[Z] = 0 \quad \text{for some } Z \in \text{ri}(\text{Sol}(\text{ESDP})) \not\Rightarrow \|x_i - x_i^{\text{true}}\| \approx 0.$$

Proof is by counterexample.

An example of sensitivity of ESDP solns to measurement noise:

Input distance data: $\epsilon > 0$

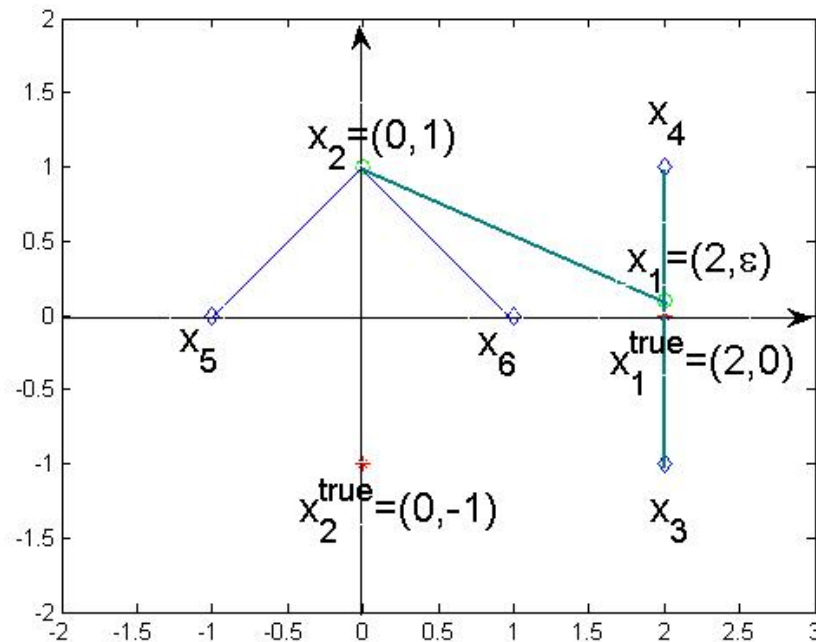
$d_{12} = \sqrt{4 + (1 - \epsilon)^2}$, $d_{13} = 1 + \epsilon$, $d_{14} = 1 - \epsilon$, $d_{25} = d_{26} = \sqrt{2}$; $m = 2$, $n = 6$.



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Thus, even when $Z \in \text{Sol}(\text{ESDP})$ is unique, $\text{tr}_i[Z] = 0$ fails to certify accuracy of x_i in the noisy case!

Robust ESDP

For each $(i, j) \in \mathcal{A}$, fix $\rho_{ij} > |\delta_{ij}|$ ($\rho > |\delta|$).

$\text{Sol}(\rho\text{ESDP})$ denotes the set of $Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix}$ satisfying

$$\begin{aligned} \begin{bmatrix} y_{ii} & y_{ij} & x_i^T \\ y_{ij} & y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \succeq 0 & \quad \forall (i, j) \in \mathcal{A}^s \\ \begin{bmatrix} y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \succeq 0 & \quad \forall i = 1, \dots, m \\ |y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2| \leq \rho_{ij} & \quad \forall (i, j) \in \mathcal{A}^a \\ |y_{ii} - 2y_{ij} + y_{jj} - d_{ij}^2| \leq \rho_{ij} & \quad \forall (i, j) \in \mathcal{A}^s. \end{aligned}$$

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Note: $Z^{\text{true}} = \begin{bmatrix} X^{\text{true}} & I \end{bmatrix}^T \begin{bmatrix} X^{\text{true}} & I \end{bmatrix} \in \text{Sol}(\rho\text{ESDP})$.

Let

$$\begin{aligned}
 Z^{\rho, \delta} &:= \arg \min_{Z \in \text{Sol}(\rho \text{ESDP})} - \sum_{(i,j) \in \mathcal{A}^s} \ln \det \left(\begin{bmatrix} y_{ii} & y_{ij} & x_i^T \\ y_{ij} & y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \right) \\
 &- \sum_{i \leq m} \ln \det \left(\begin{bmatrix} y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \right).
 \end{aligned}$$

Fact 4 (P, Tseng '09): $\exists \eta > 0, \bar{\rho} > 0$ such that for each i ,

$$\text{tr}_i(Z^{\rho, \delta}) < \eta, \text{ for some } 0 < \rho < \bar{\rho}e \implies \lim_{|\delta| < \rho \rightarrow 0} x_i^{\rho, \delta} = x_i^{\text{true}}.$$

Moreover,

$$\|x_i^{\rho, \delta} - x_i^{\text{true}}\| \leq \sqrt{2|\mathcal{A}^s| + m} (\text{tr}_i[Z^{\rho, \delta}])^{\frac{1}{2}} \quad 0 \leq |\delta| < \rho.$$

LPCGD Algorithm

Let $h_a(t) := \frac{1}{2}(t - a)_+^2 + \frac{1}{2}(-t - a)_+^2$ ($|t| \leq a \iff h_a(t) = 0$).

$$\begin{aligned}
 f_\mu(Z) := & \sum_{(i,j) \in \mathcal{A}^a} h_{\rho_{ij}}(y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2) \\
 & + \sum_{(i,j) \in \mathcal{A}^s} h_{\rho_{ij}}(y_{ii} - 2y_{ij} + y_{jj} - d_{ij}^2) \\
 & - \mu \sum_{(i,j) \in \mathcal{A}^s} \ln \det \left(\begin{bmatrix} y_{ii} & y_{ij} & x_i^T \\ y_{ij} & y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \right) \\
 & - \mu \sum_{i \leq m} \ln \det \left(\begin{bmatrix} y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \right).
 \end{aligned}$$

- f_μ is partially separable, strictly convex & diff. on its domain.
- For each $\rho > |\delta|$, $\operatorname{argmin} f_\mu \rightarrow Z^{\rho,\delta}$ as $\mu \rightarrow 0$.
- In the noiseless case ($\delta = 0$), if $\rho > 0$ is small, then $Z^{\rho,0} \approx \text{some } Z \in \operatorname{ri}(\operatorname{Sol}(\operatorname{ESDP}))$.

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Use Block Coordinate Gradient Descent

LPCGD Algorithm:

Given $Z \in \text{dom } f_\mu$, compute gradient $\nabla_{Z_i} f_\mu$ of f_μ w.r.t. $Z_i := \{x_i, y_{ii}, y_{ij} : (i, j) \in \mathcal{A}\}$ for each i .

- If $\|\nabla_{Z_i} f_\mu\| \geq \max\{\mu, 1e - 7\}$ for some i , update Z_i by moving along the Newton direction $-\left(\partial_{Z_i Z_i}^2 f_\mu\right)^{-1} \nabla_{Z_i} f_\mu$ with Armijo stepsize rule.
- Decrease μ when $\|\nabla_{Z_i} f_\mu\| < \max\{\mu, 1e - 7\}$ for all i .

$\mu_{\text{initial}} = 10$, $\mu_{\text{final}} = 1e - 14$. Decrease μ by a factor of 10.

Coded in Fortran. Computation easily distributes.

Simulation Results

- Compare ρ ESDP, as solved by LPCGD, and ESDP, as solved by Sedumi (with the interface to Sedumi coded by Wang et al.).
- Uniformly generate $\{x_1^{\text{true}}, \dots, x_n^{\text{true}}\}$ in $[-.5, .5]^2$, $m = .9n$. Two pts are neighbors iff $\text{dist} < rr$. Set

$$d_{ij} = \|x_i^{\text{true}} - x_j^{\text{true}}\| \cdot |1 + \sigma\epsilon_{ij}|,$$

where $\epsilon_{ij} \sim N(0, 1)$.

- Sensor i is judged as “accurately positioned” if

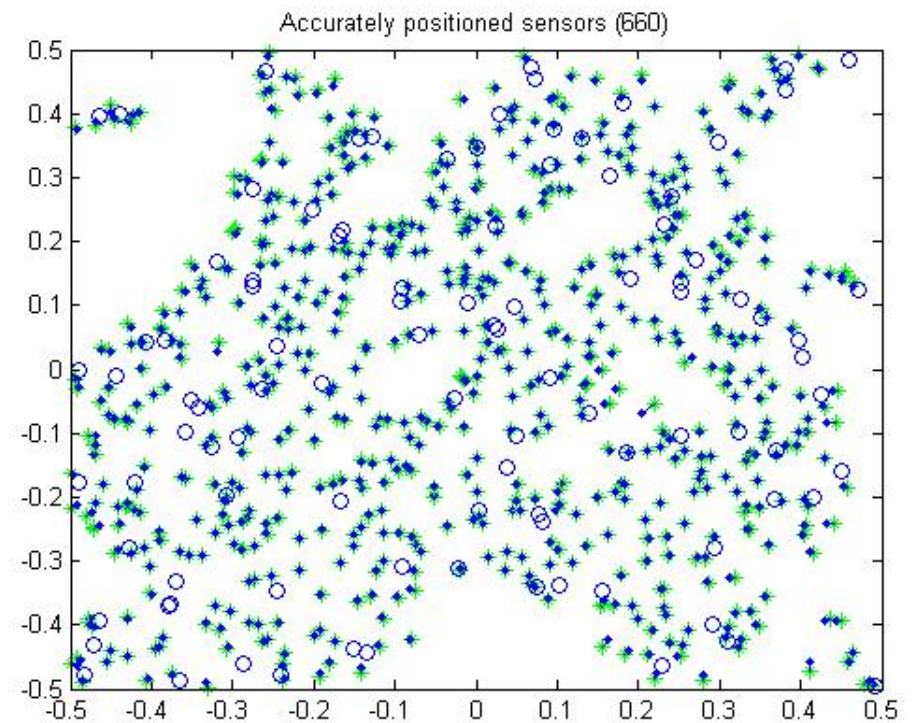
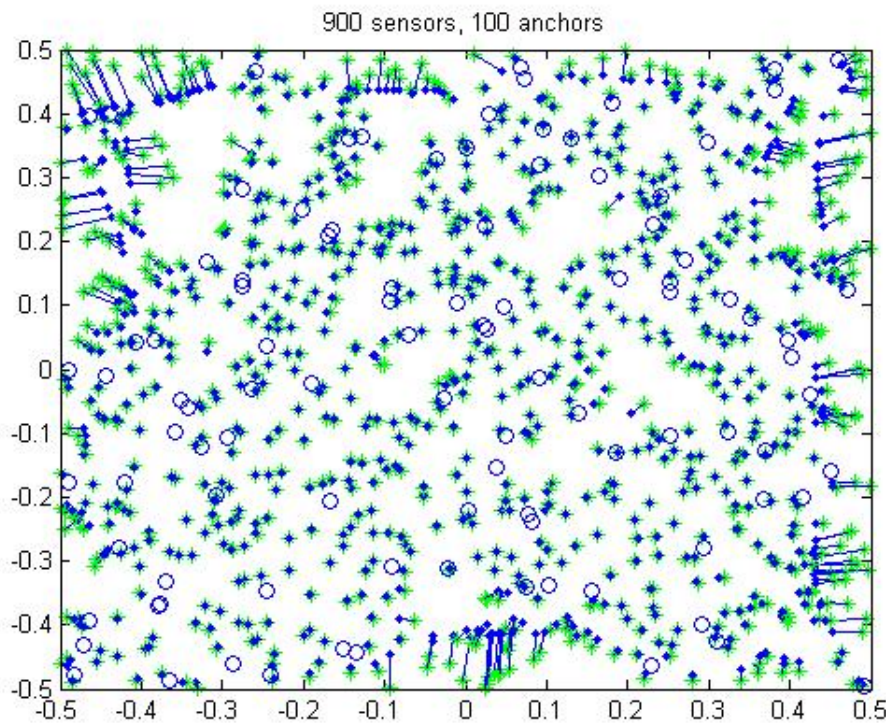
$$\text{tr}_i[Z^{\text{found}}] < (.01 + 30\sigma)\bar{d}_i^2.$$

Simulation Results

				ρ ESDP _{LPCGD}	ESDP _{Sedumi}
n	m	σ	rr	cpu / m_{ap} / err_{ap}	cpu(cpus) / m_{ap} / err_{ap}
1000	900	0	.06	7/662/1.7e-3	182(104)/669/2.1e-3
1000	900	.01	.06	5/660/2.2e-2	119(42)/720/3.1e-2
2000	1800	0	.06	26/1762/3.1e-4	1157(397)/1742/3.9e-4
2000	1800	.01	.06	20/1699/1.4e-2	966(233)/1746/2.4e-2
10000	9000	0	.02	77/7844/2.3e-3	16411(1297)/6481/2.5e-3
10000	9000	.01	.02	63/8336/1.0e-2	16368(1264)/8593/8.7e-3

- cpu(sec) times are on an HP DL360 workstation, running Linux 3.5. ESDP is solved by Sedumi; cpus:= time taken in running Sedumi.
- Take $\rho_{ij} = d_{ij}^2 \cdot ((1 - 2\tilde{\sigma})^{-2} - 1)$; $\tilde{\sigma} = \max\{\sigma, 1e - 6\}$.
- $m_{ap} := \#$ accurately positioned sensors.
 $err_{ap} := \max_{i \text{ accurate. pos.}} \|x_i^{\text{found}} - x_i^{\text{true}}\|$.

900 sensors, 100 anchors, $rr = 0.06$, $\sigma = 0.01$, solving ρ ESDP by LPCGD. x_i^{true} denoted by green asterisks, x_i^{LPCGD} denoted by blue dots, anchors denoted by circles. x_i^{true} and x_i^{LPCGD} joined by blue straight line.



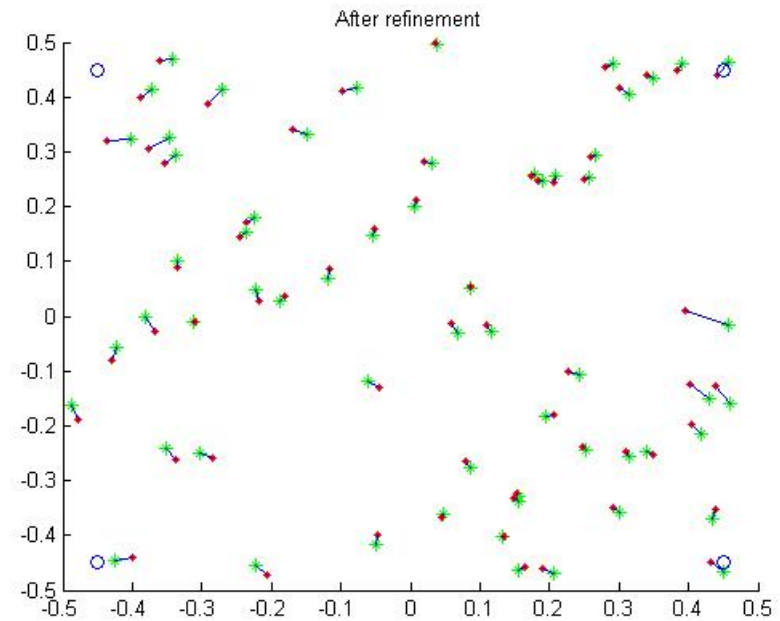
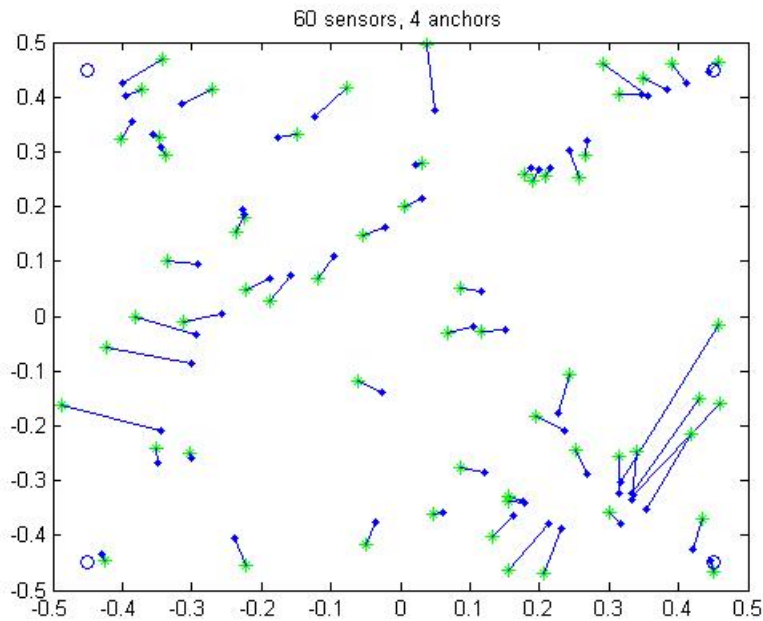
Solution Refinement

Update $x_i \leftarrow x_i - \alpha \nabla_{x_i} \hat{f}(X)$, $\forall i$, where $\hat{f}(X) := \sum_{(i,j) \in \mathcal{A}} (\|x_i - x_j\| - d_{ij})^2$.

Solution Refinement

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60 sensors, 4 anchors, $rr = 0.3$, $\sigma = 0.1$, solving ρ ESDP by LPCGD. x_i^{true} denoted by green asterisks, x_i^{LPCGD} denoted by blue dots, x_i^{refine} denoted by red dots, anchors denoted by circles. x_i^{true} and x_i^{LPCGD} / x_i^{refine} joined by blue straight line.



Conclusion & Extension

- ESDP is sensitive to dist. measurement noise. Lack soln accuracy certificate.
- ρ ESDP has soln accuracy certificate when $\rho > |\delta|$. Can $\rho > |\delta|$ be relaxed? seems not.
- ESDP/ ρ ESDP solns can be refined by performing gradient descent. RMSD is improved when noise is high ($\sim 10\%$).
- Extension of our analysis to SOS relaxation (Nie '06)? Adding range (upper bd/lower bd) constraints?

Thanks for coming! ☺